

UNIVERSITÀ DEGLI STUDI DELL'AQUILA  
**Non-Cooperative Networks: Mid-term Evaluation**  
 Wednesday, November 17th, 2021 – Prof. Guido Proietti

	Last name: .....	First name: .....	ID number: .....	Points
EXERCISE 1				
EXERCISE 2				
TOTAL				

**EXERCISE 1: Multiple-choice questions (20 points)**

**Remark:** Only one choice is correct. Use the enclosed grid to select your choice. A correct answer scores 3 points, while a wrong answer receives a  $-1$  penalization. You are allowed to omit an answer. If you wrongly select an answer, just make a circle around the wrong  $\times$  (i.e., in the following way  $\otimes$ ) and select through a  $\times$  the newly selected answer. A question collecting more than one answer will be considered as omitted. The final score will be given by summing up all the obtained points (0 for a missing answer), and then normalizing to 20.

1. Which of the following claim is false as far as the *Dominant Strategy Equilibrium* is concerned?
  - a) if  $p_i$  is a cost, it is a strategy combination  $s^* = (s_1^*, \dots, s_N^*)$ , such that for each player  $i$  and for any possible alternative strategy profile  $s = (s_1, \dots, s_i, \dots, s_N)$ ,  $p_i(s_1, \dots, s_i^*, \dots, s_N) \leq p_i(s_1, \dots, s_i, \dots, s_N)$
  - b) if  $p_i$  is a utility, it is a strategy combination  $s^* = (s_1^*, \dots, s_N^*)$ , such that for each player  $i$  and for any possible alternative strategy profile  $s = (s_1, \dots, s_i, \dots, s_N)$ ,  $p_i(s_1, \dots, s_i^*, \dots, s_N) \geq p_i(s_1, \dots, s_i, \dots, s_N)$
  - \*c) if  $p_i$  is a cost, it is a strategy combination  $s^* = (s_1^*, \dots, s_N^*)$ , such that for each player  $i$  and for any possible alternative strategy profile  $s = (s_1, \dots, s_i, \dots, s_N)$ ,  $p_i(s_1^*, \dots, s_i^*, \dots, s_N^*) \leq p_i(s_1, \dots, s_i, \dots, s_N)$
  - d) Dominant Strategy is the best possible response to any strategy of other players
2. Which of the following claim is true as far as the *Nash Equilibrium* (NE) is concerned?
  - a) It can be shown that there exist games for which finding a NE in mixed strategies is NP-hard
  - b) Any game with a finite set of players and a finite set of strategies has a NE of pure strategies
  - c) In the Head and Tail game, it does not exist a NE in mixed strategies
  - \*d) Finding a NE in pure strategies is NP-hard for many games
3. Which of the following claim is false for the Prisoner's Dilemma game:
  - a) It does admit a Nash equilibrium
  - b) It does admit a dominant strategy equilibrium
  - c) It has a Price of Anarchy equal to 5
  - \*d) It has a Price of Stability equal to  $4/3$
4. How the Price of Anarchy is defined for a game in which the social choice function  $C$  has to be minimized ( $S$  is the set of Nash equilibria)?
  - a)  $PoA = \inf_{s \in S} \frac{C(s)}{C(OPT)}$
  - b)  $PoA = \sup_{s \in S} \frac{C(OPT)}{C(s)}$
  - \*c)  $PoA = \sup_{s \in S} \frac{C(s)}{C(OPT)}$
  - d)  $PoA = \inf_{s \in S} \frac{C(OPT)}{C(s)}$
5. Which of the following claim is false in the Pigou's game:
  - a) It does admit a Nash equilibrium
  - \*b) It does not admit a dominant strategy equilibrium
  - c) The cost of the optimal flow is 0.75
  - d) The cost of the Nash flow is 1.
6. Which of the following claim is false as far as the Global Connection Game (GCG) is concerned?
  - a) A GCG is a potential game
  - \*b) The PoA of a GCG with  $k$  players is at most  $H_k$
  - c) Finding a best possible NE in a GCG is NP-hard
  - d) A best response for a player in a GCG can be found in polynomial time
7. In a Local Connection Game with  $k$  players and building cost  $\alpha \geq 0$ , which of the following claim is true?
  - a) A LCG is a potential game
  - b) for  $\alpha \leq 2$ , the complete graph is a stable solution
  - c) for  $\alpha \geq 1$ , the star is an optimal solution
  - \*d) Finding a best response for a player in a LCG is NP-hard
8. Which of the following claim is false as far as the Vickrey's Auction is concerned?
  - a) It satisfies voluntary participation
  - \*b) It does not make use of a Clarke payment scheme
  - c) It is a VCG-mechanism
  - d) It is associated with a single-parameter problem
9. Which of the following corresponds to the definition of the Ackermann function?
  - a)  $A(1, j) = 2^j$  for  $j \geq 1$ ,  $A(i, 1) = A(i - 1, 2)$  for  $i \geq 2$ ,  $A(i, j) = A(i - 1, A(i - 1, j - 1))$  for  $i, j \geq 2$
  - \*b)  $A(1, j) = 2^j$  for  $j \geq 1$ ,  $A(i, 1) = A(i - 1, 2)$  for  $i \geq 2$ ,  $A(i, j) = A(i - 1, A(i, j - 1))$  for  $i, j \geq 2$
  - c)  $A(1, j) = 2^j$  for  $j \geq 1$ ,  $A(i, 1) = A(i - 1, 2)$  for  $i \geq 2$ ,  $A(i, j) = A(i, A(i, j - 1))$  for  $i, j \geq 2$
  - d)  $A(1, j) = 2^j$  for  $j \geq 1$ ,  $A(i, 1) = A(i - 1, 2)$  for  $i \geq 2$ ,  $A(i, j) = A(1, A(i, j - 1))$  for  $i, j \geq 2$
10. In the one-parameter mechanism for the single-source shortest path tree problem, which payment will receive an edge  $e$  belonging to the solution?
  - a)  $p_e = r_e w_e(r) + \int_0^\infty w_e(r - e, z) dz$
  - \*b)  $p_e = r_e w_e(r) + \int_{r_e}^\infty w_e(r - e, z) dz$
  - c)  $p_e = -r_e w_e(r) + \int_{r_e}^\infty w_e(r - e, z) dz$
  - d)  $p_e = r_e w_e(r) + \int_0^{r_e} w_e(r - e, z) dz$

**Answer Grid**

	Question									
Choice	1	2	3	4	5	6	7	8	9	10
a										
b										
c										
d										

**EXERCISE 2: Open question (10 points)**

**Remark:** Select at your choice one out of the following two questions, and address it exhaustively.

1. Describe and analyze the global connection game.
2. Describe and analyze the one-parameter mechanism for the private-edge shortest path tree problem.