

**Laurea Magistrale in Informatica**  
**Formal Methods - Rewriting (2010-2011)**

**Intermediate Written Exam**

**December 2nd, 2010**

1. Given the signature  $\Sigma = \{a, f, g, h\}$  and variables  $x, y, x', y' \in V$ , compute the most general unifier (if it exists) of the following pairs of terms:

i)  $t_1 = g(h(x), g(x, y))$  and  $t_2 = g(x', g(x', x'))$ ;

ii)  $t_1 = f(x, h(y), h(x))$  and  $t_2 = f(a, x', h(y'))$ ;

iii)  $t_1 = f(x, g(x, y), h(y))$  and  $t_2 = f(h(x'), g(h(y'), y'), h(a))$ .

2. Given the terms  $t_1 = f(g(x, y), z)$  and  $t_2 = f(f(x', y'), z')$ , say if  $t_1$  can be syntactically unified with  $t_2$  and subterms of  $t_2$ , and give the most general unifiers (if they exist).

3. Consider the following TRS  $R$  on the signature  $\Sigma = \{a, f, g, h, k\}$ :

$$\begin{aligned} g(a, x) &\rightarrow a \\ g(h(x), y) &\rightarrow f(k(y), g(x, y)) \\ k(a) &\rightarrow a \end{aligned}$$

i) Give a reduction ordering on terms such that  $R$  is terminating with respect to such an ordering. Show the formal steps that justify your answer.

ii) Reduce the term  $t = g(h(h(a)), a)$  to normal form in  $R$ , by applying all the possible reduction steps from  $t$  in  $R$  and showing the rule applied, the position of the redex and the matching substitution for each reduction step.

4. Consider the following TRS  $R$  on the signature  $\Sigma = \{a, b, f, g, k\}$ :

$$\begin{aligned} k(a) &\rightarrow b \\ k(f(x)) &\rightarrow x \\ g(x, b) &\rightarrow x \\ g(x, a) &\rightarrow f(x) \\ g(x, f(y)) &\rightarrow f(g(x, y)) \end{aligned}$$

Give a reduction ordering on terms such that  $R$  is terminating with respect to such an ordering. Show the formal steps that justify your answer.

5. Given an ARS  $\mathcal{A} = \langle A, \longrightarrow \rangle$ , the *unique normal form* property (UN for short) can be defined as follows: for all  $a, b \in A$  if  $a \xrightarrow{*} b$  and both  $a$  and  $b$  are normal forms, then  $a \equiv b$ . Prove that if an ARS  $\mathcal{A}$  is UN and WN, then  $\mathcal{A}$  is confluent.