## Algorithm Design Laboratory with Applications

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## **Problem:** Bookshelf.

You are shopping for a new bookshelf to organize your collection of  $n \in \mathbb{N}^+$  books indexed with the integers in  $\{1, \ldots, n\}$ . To fit in your room, the bookshelf can only be tall enough to accommodate 3 shelves, that is, 3 rows of books.

The *i*-th book in your collection (i = 1, ..., n) has thickness  $t_i \in \mathbb{N}^+$ , therefore a bookshelf of width W will be able to fit all the books if and only if there is a partition  $\mathcal{B} = \{B_1, B_2, B_3\}$  of  $\{1, ..., n\}$  into 3 sets  $B_1, B_2, B_3$  (corresponding to the books to place on the top, middle, and bottom shelf, respectively) such that  $\forall B \in \mathcal{B}, \sum_{i \in B} t_i \leq W$ .

Design an algorithm that, given n and  $t_1, \ldots, t_n$ , computes the minimum width W of a bookshelf capable of fitting all books.

**Input.** The input consists of a set of instances, or *test-cases*, of the previous problem. The first line contains the number T of test-cases. Each test case consists of a single line containing the integer n followed by the values  $t_1, \ldots, t_n$ .

**Output.** The output consists of T lines. The *i*-th line is the answer to the *i*-th test-case and contains the minimum width W of a bookshelf capable of fitting all books.

Assumptions.  $1 \le T \le 10$ ;  $1 \le n \le 2^6$ ;  $\forall i = 1, \dots, n, 1 \le t_i \le 2^6$ . Example.



Input (corresponding to the above example, book labels represent their thickness):

1 17 2 1 3 1 1 2 4 2 1 1 4 2 1 1 1 2 2

*Output (a possible arrangement is shown in the figure above):* 

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**Requirements.** Your algorithm should require time  $O(n\tau^2)$ , where  $\tau = \sum_{i=1}^{n} t_i$  (with reasonable hidden constants).

Notes. A reasonable implementation should not require more than 3 seconds for each input file.