

# Algorithm Design Laboratory with Applications

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**Problem:** *A problem of two trees.*

A computer scientist with green fingers wants to remove a big tree<sup>1</sup> from his garden and has hired a moving company to transport it away. The moving company owns a truck that can carry a maximum weight of  $W \in \mathbb{N}^+$ .

The tree is too heavy to be transported at once, but it can be cut into smaller (and lighter) pieces. The computer scientist has modelled this problem with a rooted tree<sup>2</sup>  $T = (V, E)$  in which each edge  $e \in E$  has a weight  $w(e) \in \mathbb{N}^+$ . The weight  $w(T')$  of a subtree  $T'$  of  $T$  is the sum of the weights  $w(e)$  of all edges  $e$  in  $T'$ .

Cutting a tree  $T$  in one of its internal vertices  $v$  means splitting  $T$  into two  $1 + c(v)$  trees  $T_0, T_1, \dots, T_{c(v)}$ , where  $c(v)$  is the number of children  $u_1, \dots, u_{c(v)}$  of  $v$  in  $T$ . In details:

- $T_0$  is unique tree containing  $v$  in the forest obtained by deleting  $u_1, \dots, u_{c(v)}$  from  $T$ .  $T_0$  is rooted in the same root as  $T$ .
- For  $i = 1, \dots, c(v)$ ,  $T_i$  is the subtree of  $T$  induced by  $v$  and all the descendants of  $u_i$  in  $T$ .  $T_i$  is rooted in  $v$ .

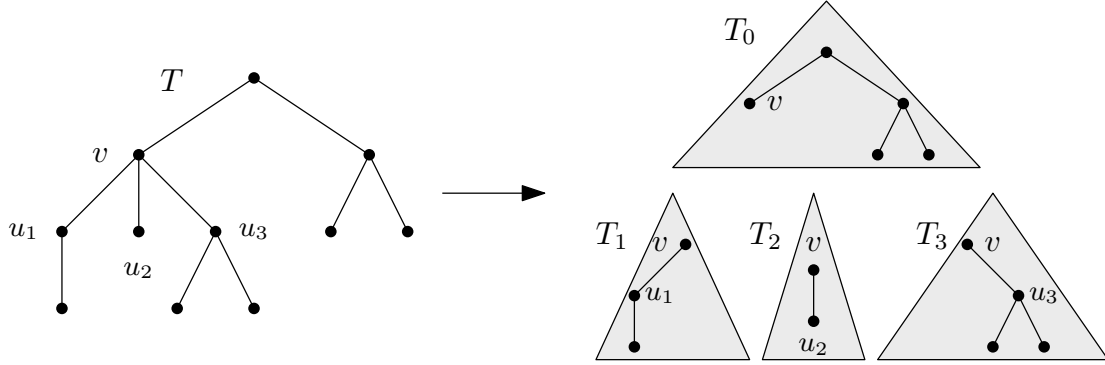


Figure 1: An example of the trees  $T_0, \dots, T_{c(v)}$  resulting from cutting  $T$  in  $v$ . Edge weights are not shown.

Help the computer scientist design a fast algorithm that determines the minimum number  $\eta(T, W)$  of cuts needed to decompose  $T$  into a forest  $F$  in which each tree  $T' \in F$  has a weight  $w(T')$  of at most  $W$ .

**Input.** The input consists of a set of instances, or *test-cases*, of the previous problem. The first line of the input contains the number  $C$  of test-cases. Each test-case is described by 3 lines. The first line of each test-case contains  $W$  and the number  $n > 0$  of vertices of  $T$ . The vertices of  $T$  are indexed from 0 to  $n - 1$ , and the root of  $T$  is the vertex with index 0. The second line contains  $n - 1$  integers  $p_1, \dots, p_{n-1}$  separated by a space, where  $p_i$  is the index of the parent (in  $T$ ) of the unique vertex with index  $i$ . The third and final line contains  $n - 1$  integers. The  $i$ -th of these integers is the weight  $w(e)$  of the edge  $e = (v, u)$  connecting the vertex  $u$  with index  $i$  to its parent  $v$  in  $T$ .

**Output.** The output consists of  $C$  lines. The  $i$ -th line is the answer to the  $i$ -th test-case and contains the integer  $\eta$ .

**Assumptions.**  $1 \leq C \leq 10$ ;  $1 \leq n < 2^{18}$ ;  $\forall e \in E, 1 \leq w(e) \leq 2^{10}$ ;  $\max_{e \in E} w(e) \leq W < 2^{31}$ .

<sup>1</sup>Unlike the trees the computer scientist is used to, this tree is made of solid wood and its roots are at the bottom.

<sup>2</sup>The kind of tree the computer scientist is familiar with.

Vertices of  $T$  with the same parent have consecutive indices in the input.

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1
10 21
0 0 1 1 1 2 2 3 3 5 7 7 8 8 10 10 10 11 11 11
3 2 2 1 5 1 1 2 3 1 2 2 3 2 2 5 1 1 2 1

```

4

**Notes.** A reasonable implementation should not require more than 0.5 seconds for each input file.